

DĂNUȚ M A R C U (\*)

## Note on the spanning trees of a connected graph (\*\*)

### Introduction

This paragraph is meant to present some definitions and results that are necessary to follow the further notes, our graph theoretic terminology being fairly standard [1], [2].

A *spanning tree* of a graph (an undirected graph) is a tree of the graph that contains all the vertices in the graph. It is well-known that a graph is connected if and only if contains a spanning tree [1].

A *cocycle* is a minimal set of edges in a graph the removal of which will increase the number of connected components in the remaining subgraph, whereas the removal of any its proper subset will not. It follows that in a connected graph the removal of a cocycle will separate the graph into two parts. This suggests an alternative way of defining a cocycle. Let the vertices in a connected component of a graph be divided into two subsets such that every two vertices in one subset are connected by a chain that contains only vertices in the subset. Then, the set of edges joining the vertices in the two subsets is a cocycle.

Since the removal of any edge  $e$  from a spanning tree  $T$  breaks the spanning tree up into two trees (that may consist of a single vertex), it follows that for every edge in a spanning tree there is a unique corresponding cocycle of the graph, called *fundamental cocycle* with respect to  $e$  and  $T$ . A graph is said to be obtained from  $G$  by *open-circuiting* and edge  $e$ , if it is obtained removing  $e$  from  $G$ . The open-circuiting of more than one edge is defined similarly (e.g., see [3]<sub>2</sub>).

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(\*) Indirizzo: Faculty of Mathematics, University of Bucharest, Academiei 14, 70109 Bucharest, Romania.

(\*\*) Ricevuto: 28-VI-1983.

A graph (possibly a multigraph) is said to be obtained from  $G$  by *short-circuiting* an edge  $e$ , if it is obtained by identifying, i.e., combining into one vertex, the two end-vertices of  $e$  and then removing  $e$ . The short-circuiting of more than one edge is defined similarly (e.g., see [3]<sub>2</sub>).

If  $T = \{e_1, e_2, \dots, e_m\}$  is a spanning tree of a connected graph  $G$  (according to [1],  $G$  contains  $m + 1$  vertices), we shall denote by  $G(T, e_{i_1}, e_{i_2}, \dots, e_{i_k})$ ,  $k = 2, 3, \dots, m$  the graph (possibly the multigraph) obtained from  $G$  by open-circuiting the edges  $e_{i_1}, e_{i_2}, \dots, e_{i_{k-1}}$  and short-circuiting the edges  $\{f_1, f_2, \dots, f_t\} = T - \{e_{i_1}, e_{i_2}, \dots, e_{i_k}\}$ ,  $t = m - k$ .

Remark. It is easy to check that the order, in which the above two graph's transformations are applied to obtain  $G(T, e_{i_1}, e_{i_2}, \dots, e_{i_k})$ , can be arbitrary.

The main results. Let  $G = (V, E)$  be a connected graph ( $V$  is the set of vertices, and  $E$  the set of edges),  $T$  a spanning tree of  $G$  and  $e \in T$ . We shall denote by  $C(e, T)$  the fundamental cocycle with respect to  $e$  and  $T$ .

Let  $v \in V$  and  $e \in E$ , arbitrary chosen, such that  $e$  is incident with the vertex  $v$ . We denote by  $C(v)$  the cocycle associated to the bipartition  $(\{v\}, V - \{v\})$ . Let  $A(e)$  the set of spanning trees that contain the edge  $e$  and  $\bar{A}(e)$  the set of spanning trees that do not contain  $e$ .

Lemma ([3]<sub>1</sub>, Theorem 3).

$$\bar{A}(e) = \bigcup_{\substack{T \in A(e) \\ b \in C(e, T) \cap C(v) \\ b \neq e}} \{(T - \{e\}) \cup \{b\}\}.$$

Let now  $T = \{e_1, e_2, \dots, e_m\}$  be a fixed spanning tree of a connected graph  $G = (V, E)$ , such that for every  $k = 1, 2, \dots, m$  the subset  $\{e_1, e_2, \dots, e_k\}$  forms a connected subgraph of  $G$  (obviously, the edges of a spanning tree can be always arranged in this way). Replacing the edges  $e_{i_1}, e_{i_2}, \dots, e_{i_k}$  ( $k \leq m$ ) of  $T$  with  $k$  distinct edges from  $E - T$ , we obtain a spanning subgraph  $S$  of  $G$ . If  $S$  is connected, then it is a spanning tree of  $G$ . So, let us denote by  $A(T, e_{i_1}, e_{i_2}, \dots, e_{i_k})$  the set of the distinct spanning trees of  $G$  obtained by replacing, in all the possible ways, the edges  $e_{i_1}, e_{i_2}, \dots, e_{i_k}$  of  $T$  with  $k$  distinct edges from  $E - T$  (obviously,  $A(T, e_{i_1}, e_{i_2}, \dots, e_{i_k})$  can be empty).

Theorem. If  $1 \leq i_1 < i_2 < \dots < i_k \leq m$ , then

$$A(T, e_{i_1}, e_{i_2}, \dots, e_{i_k}) = \bigcup_{\substack{T' \in A(T, e_{i_1}, e_{i_2}, \dots, e_{i_{k-1}}) \\ b \in C(e_{i_k}, T) \cap C(e_{i_k}, T') \\ b \neq e_{i_k}}} \{(T' - \{e_{i_k}\}) \cup \{b\}\}.$$

Proof. Let  $\tilde{G} = G(T, e_{i_1}, e_{i_2}, \dots, e_{i_k})$  and  $T' \in A(T, e_{i_1}, e_{i_2}, \dots, e_{i_{k-1}})$  arbitrary chosen. We denote by  $\tilde{T}'$  what it remains from  $T'$  in  $\tilde{G}$ , i.e.,  $\tilde{T}' = T' - \{f_1, f_2, \dots, f_i\}$ . Obviously, since  $T'$  does not contain the edges  $e_{i_1}, e_{i_2}, \dots, e_{i_{k-1}}$ ,  $\tilde{T}'$  is a spanning tree of  $\tilde{G}$  (thus  $\tilde{G}$  is connected), and  $\tilde{T}' \in \tilde{A}(e_{i_k})$ .

Because the set  $\{e_1, e_2, \dots, e_{i_k}\}$  forms a connected subgraph of  $T$  and  $\{e_{i_1}, e_{i_2}, \dots, e_{i_k}\} \subseteq \{e_1, e_2, \dots, e_{i_k}\}$ , it results that one of the two connected components of  $T - \{e_{i_k}\}$  is either an isolated vertex (if  $e_{i_k}$  is incident with a pendant vertex of  $T$ ) or a subgraph whose set of edges is contained in  $T - \{e_{i_1}, e_{i_2}, \dots, e_{i_k}\}$ .

Let  $v$  be the vertex of this connected component of  $T - \{e_{i_k}\}$  that is incident with  $e_{i_k}$  (this component is reduced to  $v$  in  $\tilde{G}$ ). Obviously, the procedure for obtaining  $\tilde{G}$  does not affect the fundamental cocycle  $C(e_{i_k}, T)$  that coincides with  $\tilde{C}(v)$  in  $\tilde{G}$ . On the other hand, the fundamental cocycles  $\tilde{C}(e_{i_k}, \tilde{T}')$  and  $C(e_{i_k}, T')$  differ one of another only by edges from the set  $\{e_{i_1}, e_{i_2}, \dots, e_{i_{k-1}}\}$ , these edges being open-circuited. Thus, the following equality holds

$$(1) \quad \tilde{C}(v) \cap \tilde{C}(e_{i_k}, \tilde{T}') = C(e_{i_k}, T) \cap C(e_{i_k}, T').$$

From (1) and Lemma (used for the connected graph  $\tilde{G}$ ) it follows that

$$(2) \quad \tilde{A}(e_{i_k}) = \bigcup_{\substack{\tilde{T}' \in \tilde{A}(e_{i_k}) \\ b \in C(e_{i_k}, T) \cap C(e_{i_k}, T') \\ b \neq e_{i_k}}} \{(\tilde{T}' - \{e_{i_k}\}) \cup \{b\}\}.$$

According to [3]<sub>1</sub> (Theorem 4),  $\tilde{T}'$  is uniquely obtained from  $T'$  by the construction of  $\tilde{G}$ . Since  $\{f_1, f_2, \dots, f_i\} \subseteq T'$  for each  $T' \in A(T, e_{i_1}, e_{i_2}, \dots, e_{i_{k-1}})$ , then, by [3]<sub>2</sub> (theorems 5.1 and 4.1),  $T' := \tilde{T}' \cup \{f_1, f_2, \dots, f_i\}$  is unique and belongs to  $A(T, e_{i_1}, e_{i_2}, \dots, e_{i_{k-1}})$  for every  $\tilde{T}' \in \tilde{A}(e_{i_k})$ .

Thus, the sets  $\tilde{A}(e_{i_k})$  and  $A(T, e_{i_1}, e_{i_2}, \dots, e_{i_{k-1}})$  are in a bijective correspondence. Similarly, the sets  $\tilde{A}(e_{i_k})$  and  $A(T, e_{i_1}, e_{i_2}, \dots, e_{i_k})$ . Hence, the theorem is proved according to (2).

Corollary. If  $\{e_{i_1}, e_{i_2}, \dots, e_{i_k}\}$  and  $\{e_{j_1}, e_{j_2}, \dots, e_{j_r}\}$  are distinct subsets of  $T$ ,  $k, r = 1, 2, \dots, m$ , then  $A(T, e_{i_1}, e_{i_2}, \dots, e_{i_k}) \cap A(T, e_{j_1}, e_{j_2}, \dots, e_{j_r}) = \emptyset$ .

Proof. It readily follows from theorem and the definition of a fundamental cocycle.

Remark. If  $T$  is a spanning tree of  $G$  an  $e$  and edge belonging to  $T$ , then a particular case of the above theorem is the following relation

$$A(T, e) = \bigcup_{\substack{b \in C(e, T) \\ b \neq e}} \{(T - \{e\}) \cup \{b\}\}.$$

Example. Let us consider the connected graph  $G = (V, E)$  with  $V = \{v_1, v_2, v_3, v_4\}$ ,  $E = \{a = (v_1, v_2), b = (v_1, v_4), c = (v_1, v_3), d = (v_2, v_3), e = (v_2, v_4), f = (v_3, v_4)\}$  and  $T = \{a, b, d\}$  a spanning tree of  $G$  for which  $\{a\}$ ,  $\{a, b\}$  and  $\{a, b, d\}$  are connected subgraphs of  $G$ . We have

$$C(a, T) = \{a, c, e, f\}, \quad C(b, T) = \{b, e, f\}.$$

According to the above remark we obtain

$$A(T, a) = \{\{b, c, d\}, \{b, d, e\}, \{b, d, f\}\}.$$

On the other hand we have

$$C(b, \{b, c, d\}) = \{b, c, f\}, \quad C(b, \{b, d, f\}) = \{a, b, e\}, \quad C(b, \{b, d, f\}) = \{a, b, e\},$$

$$C(b, \{b, c, d\}) \cap C(b, T) = \{b, e, f\}, \quad C(b, \{b, d, e\}) \cap C(b, T) = \{b\},$$

$$C(b, \{b, d, f\}) \cap C(b, T) = \{b\},$$

and from the above theorem we obtain

$$A(T, a, b) = \{\{c, d, e\}, \{c, d, f\}\}.$$

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### References

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### Résumé

*Une théorème, concernant les arbres partiels d'un graphe connexe, est donnée.*

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